

Saving Resources on Wireless Uplinks: Models of Queue-aware Scheduling

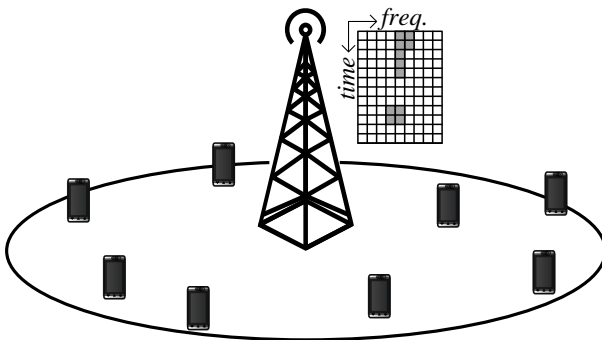


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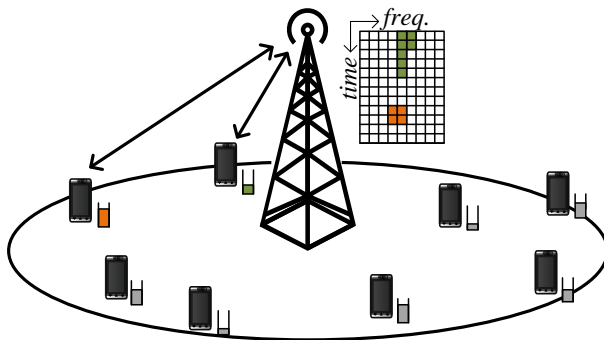
Amr Rizk
TU Darmstadt

- joint work with Markus Fidler

Cellular Uplink Scheduling



Cellular Uplink Scheduling



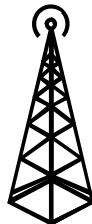
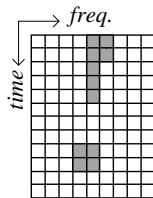
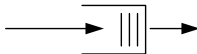
Adaptive Resource Allocation in Cellular Uplink Direction

Input metrics (LTE)

- ▶ Buffer status reports (BSR)
- ▶ Channel quality indicators (CQI)

Goal

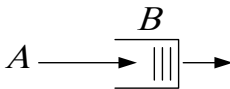
- ▶ Statistical QoS guarantee



Adaptive Resource Allocation in Cellular Uplink Direction

Input metrics (LTE)

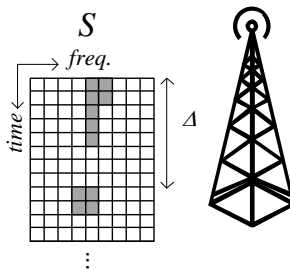
- ▶ Buffer status reports (BSR)
- ▶ Channel quality indicators (CQI)



- ▶ Arrival traffic A
- ▶ Buffer filling B

Goal

- ▶ Statistical QoS guarantee

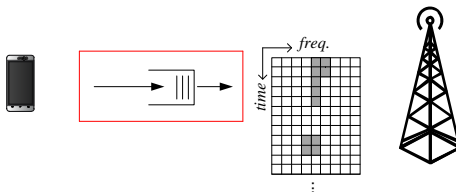


- ▶ Service S
- ▶ Scheduling epoch Δ

- ▶ Exact analysis for Poisson traffic

- ▶ Analytical framework for general arrival and service processes

Exact Analysis for Poisson Traffic

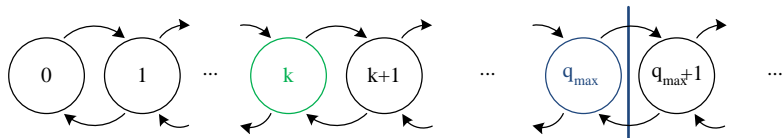




Model

- ▶ Single M/M/1 queue: fixed λ , variable $\mu(t)$
- ▶ Given λ and the queue length at epoch start
- ▶ Epoch based resource allocation

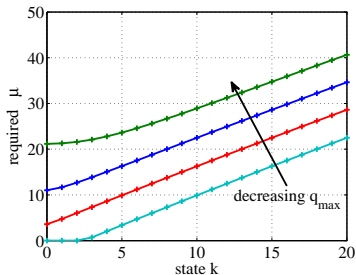
⇒ Find $\mu(t)$ that provides a probabilistic bound on the queue length at the end of the epoch



1. Queue initially in state k
2. Fix $\mu(t)$ during Δ such that
3. Probability that the queue at time Δ is longer than q_{\max} is less than ε
 - ▶ Based on the transient behavior of the $M/M/1$ queue [Kleinrock].

Model parameters: $\lambda, q_{\max}, \varepsilon, \Delta$

- ▶ Required μ for various parameters

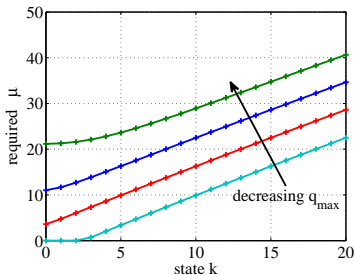


(a) $q_{\max} \in \{5, 10, 15, 20\}$

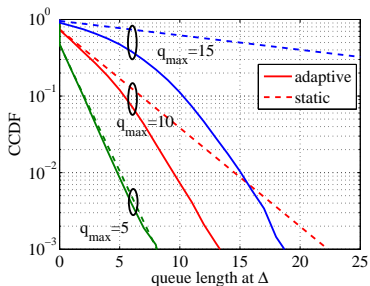
Parameters: $\lambda = 10$, $\varepsilon = 10^{-2}$, $\Delta = 1$ and 10^4 epochs for (b).

Exact Analysis for Poisson Traffic

- ▶ Required μ for various parameters
- ▶ Improvement w.r.t. the static system with equivalent $\bar{\mu}$



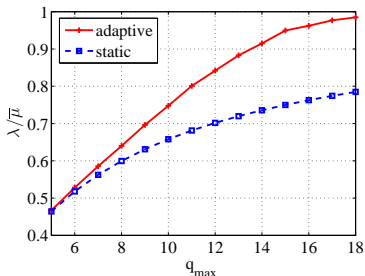
(a) $q_{\max} \in \{5, 10, 15, 20\}$



(b) Static system parameterized with $\bar{\mu}$.

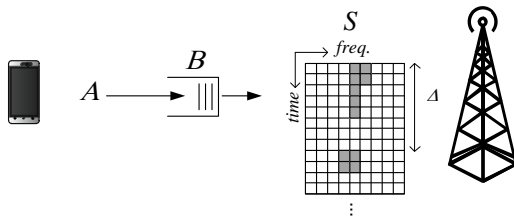
Parameters: $\lambda = 10$, $\varepsilon = 10^{-2}$, $\Delta = 1$ and 10^4 epochs for (b).

Utilization comparison:



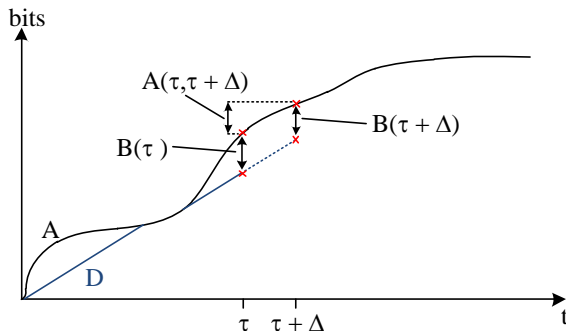
- ▶ Key relation of $\lambda\Delta$ to q_{\max} for a given ε
- ▶ Initial queue length is less helpful if the unknown traffic amount in during the epoch, i.e., $\lambda\Delta$, predominates q_{\max}
 - Operation of queue-aware scheduling is non-trivial
- ▶ Resource savings in the adaptive case → Proof of concept

Generalization w.r.t. service and arrival traffic models:



Framework: Stochastic Network Calculus

- ▶ Cumulative arrivals $A(\tau)$ resp. departures $D(\tau)$ up to time τ
- ▶ Backlog at τ : $B(\tau) = A(\tau) - D(\tau)$
- ▶ Service in $(\tau, t]$ as random process $S(\tau, t)$
- ▶ Assume strict service resp. adaptive service curve [Burchard et. al'06]



- ▶ Evaluation requires a lower bound on the service process

$$P[S(u, t) \geq S(t - u), \forall u \in [\tau, t]] \geq 1 - \varepsilon_s$$

- ▶ To derive a lower bound on the departures $D(\tau + \Delta)$

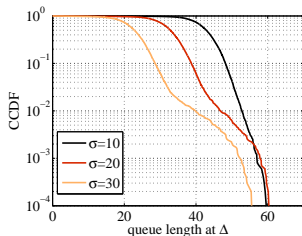
Basic block fading model for a wireless transmission [Fidler, Al-Zubaidy]

- ▶ Time slotted model with iid increments $c_i = \beta \ln(1 + \gamma_i)$
- ▶ Rayleigh fading channel: γ_i is exp distributed with parameter η
- ▶ Lower bounding function for the service process

$$S(t) = \frac{1}{\theta} \left(\ln(\varepsilon_p) - t \left[\eta + \theta \beta \ln(\eta) + \ln(\Gamma(1 - \theta \beta, \eta)) \right] \right)$$

with $\theta > 0$, incompl. Gamma fct. Γ and a violation probability $\varepsilon_s = (t - \tau)\varepsilon_p$.

Derivation using Boole's inequality, Chernoff's bound and the Laplace transform of the increments.



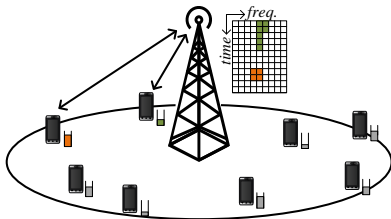
- ▶ Bound on backlog at the end of epoch $P\left[B(\tau + \Delta) \leq b_{\max}\right] \geq 1 - \varepsilon_s$
- ▶ Requirements on $\mathcal{S}(t)$: Allocate resources β during epoch Δ such that the following holds

$$\mathcal{S}(\Delta) \geq B(\tau) + A(\tau, \tau + \Delta) - b_{\max}, \quad \text{and} \quad (1)$$

$$\mathcal{S}(\tau + \Delta - u) \geq A(u, \tau + \Delta) - b_{\max}, \quad \forall u \in [\tau, \tau + \Delta] \quad (2)$$

Scenario:

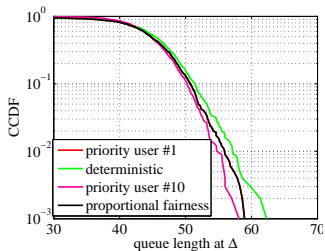
- ▶ M homogeneous, statistically independent MS channels
- ▶ Base station decides on amount of resource blocks β_j for MS $j \in [1, M]$ based on the infrequent adaptation technique
- ▶ Overall amount of resource blocks β_s in epoch Δ



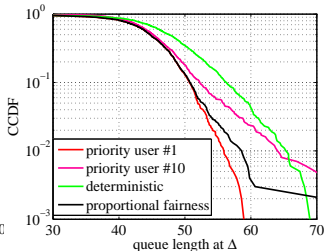
Scenario:

- ▶ M homogeneous, statistically independent MS
- ▶ Base station decides on amount of resource blocks β_j for MS $j \in [1, M]$ based on the infrequent adaptation technique
- ▶ Overall amount of resource blocks β_s in epoch Δ
- ▶ Three basic resource allocation algorithms:
 1. “deterministic fair”: j th MS receives $\hat{\beta}_j = \min\{\beta_j, \beta_s/M\}$
 2. “priority”: MS in class j receives $\hat{\beta}_j = \min\{\beta_j, \beta_s - \sum_{k=1}^{j-1} \hat{\beta}_k\}$
 3. “proportional fair emulation”: Priority scheduler with priorities reordered every epoch Δ according to a score $S_j(\tau, \tau + \Delta)/(D_j(\tau)/\tau)$ similar to [Kelly, et al. '98].

Multi-user Scheduling with Infrequent Adaptation: Simulation



(b) 90% utilization



(c) 95% utilization

- ▶ Adaptive system retains statistical backlog bound depending on scheduling algorithm
- ▶ Notable difference only at very high utilization

Parameters: $M = 10$, $\lambda = 0.65$, $\varepsilon_s = 10^{-2}$, $\Delta = 100$ slots, $b_{\max} = 65$, SNR $1/\eta = 3$ dB

Key Takeaway Points

- ▶ Poisson case: Analytical results to quantify best-case resource savings.
- ▶ Model reveals important relation of $\lambda\Delta$ to q_{\max} .
- ▶ Analytical framework identifies two regimes, one where adaptive scheduling is effective and one where it is not.
- ▶ A mathematical treatment of queue-aware scheduling that is applicable to a broad class of arrival and service processes.

1. Optimization of queueing service policies
2. Optimization of power and rate allocation in cellular systems

Difference to 1:

- ▶ online, epoch-based technique for general arrival and service processes

Difference to 2:

- ▶ dynamic programming to minimize a cost function of weighted power and rate consumption
- ▶ sample path as input